

The Semi-Analytic Landau Bootstrap

FROM ANALYTICITY TO PHENOMENOLOGY

Princeton, NJ, January 12, 2026

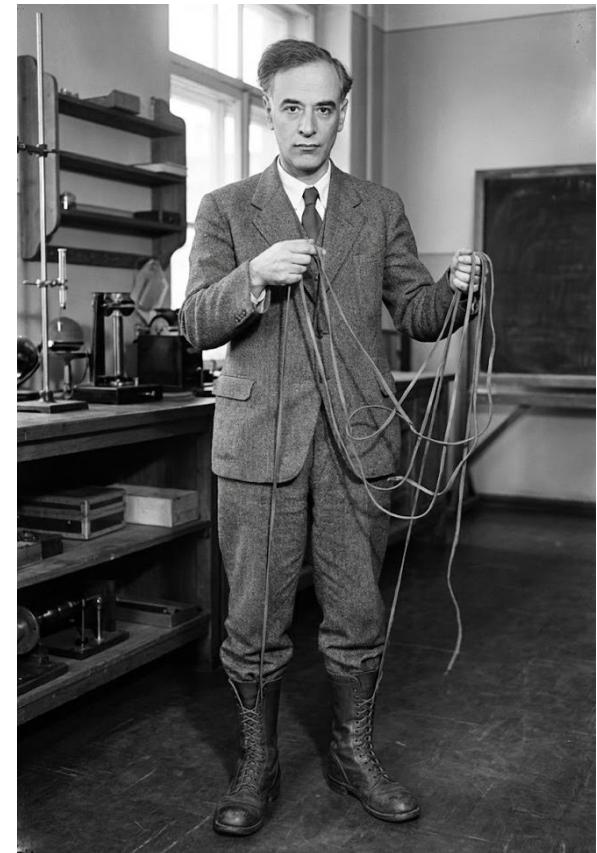
Matthew Schwartz
Harvard University



Institute for Artificial Intelligence
and Fundamental Interactions (IAIFI)

Based on

- H. Hannesdottir, A. McLeod, MDS, C. Vergu:
"Applications of the Landau Bootstrap" [2410.02424](https://arxiv.org/abs/2410.02424)
"Constraints on sequential discontinuities from the geometry of on-shell spaces" [2211.07633](https://arxiv.org/abs/2211.07633)
- O. Barrera, A. Dersy, R. Husain, MDS and X. Zhang:
"Analytic Regression of Feynman Integrals from High-Precision Numerical Sampling", [2507.17815](https://arxiv.org/abs/2507.17815)



The S-matrix bootstrap

Is the S matrix completely fixed by physical constraints?

- 1960s: best hope for strong interactions
- 1970s: Quantum Field Theory explained strong interactions
→ S matrix program on hold for 40 years

Three Modern Bootstraps

1. Non-perturbative bootstrap

- e.g. pion 2->2 scattering at low energy
- Similar to 1960's bootstrap
 - Impose unitarity, analyticity, crossing
- Use modern methods
(Hamiltonian truncation, machine learning, etc)

2. Perturbative amplitude bootstrap

- constrain full amplitude (sum of all diagrams)
- Impose symmetries and physical limits
 - dual conformal invariance, Regge limit
- 6-point in N=4 SYM at 7 loops computed this way

[Caron-Huot et al 1903.10890]

today's
talk

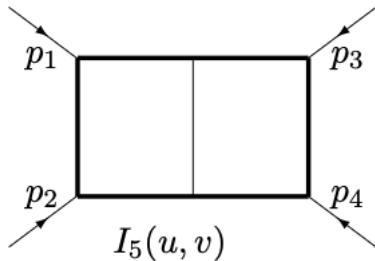
3. Feynman diagram bootstrap

- Can we compute Feynman diagrams without actually doing the integrals?
- Many examples to work with
 - Uses elements of 1 and 2

Feynman-diagram bootstrap

Q: Can we bootstrap Feynman integrals?

Rules: Don't integrate!



$$I_5(u, v) = \int \frac{d^D k_1}{(2\pi)^D} \int \frac{d^D k_2}{(2\pi)^D} \frac{1}{[k_1^2 + m^2] [k_2^2 + m^2] [(k_1 - p_1)^2 + m^2] [(k_2 - p_3)^2 + m^2] [(k_1 - k_2)^2]}$$

- Two-loop outer mass double box
- Caron-Huot and Henn 1404.2922

$$\begin{aligned}
 g_{10} = & G_{-1,0} H_{-1,-1} + G_{-1,\frac{1}{2+1}} H_{-1,-1} - G_{-1,-\frac{1+1}{2+1}} H_{-1,-1} + 2G_{0,-1} H_{-1,-1} + G_{0,0} H_{-1,-1} \\
 - & 2G_{0,\frac{1}{2}} H_{-1,-1} - G_{0,\frac{1+1}{2+1}} H_{-1,-1} - G_{0,-\frac{1+1}{2+1}} H_{-1,-1} - G_{1,0} H_{-1,-1} - G_{1,\frac{1}{2+1}} H_{-1,-1} \\
 + & G_{1,-\frac{1+1}{2+1}} H_{-1,-1} + G_{-1,0} H_{-1,0} - G_{-1,\frac{1}{2+1}} H_{-1,0} - G_{-1,-\frac{1+1}{2+1}} H_{-1,0} + G_{0,\frac{1}{2+1}} H_{-1,0} \\
 - & G_{0,-\frac{1+1}{2+1}} H_{-1,0} + G_{1,0} H_{-1,0} + G_{1,\frac{1}{2+1}} H_{-1,0} + G_{1,-\frac{1+1}{2+1}} H_{-1,0} + G_{-1,0} H_{-1,1} + G_{-1,\frac{1}{2+1}} H_{-1,1} \\
 - & G_{-1,-\frac{1+1}{2+1}} H_{-1,1} + 2G_{0,-1} H_{-1,1} + G_{0,0} H_{-1,1} - 2G_{0,\frac{1}{2}} H_{-1,1} - G_{0,-\frac{1+1}{2+1}} H_{-1,1} \\
 - & G_{1,0} H_{-1,1} - G_{1,\frac{1}{2+1}} H_{-1,1} + G_{1,-\frac{1+1}{2+1}} H_{-1,1} - G_{-1,0} H_{-1,1} + G_{-1,-\frac{1+1}{2+1}} H_{-1,1} \\
 - & G_{0,\frac{1}{2+1}} H_{0,-1} + G_{0,-\frac{1+1}{2+1}} H_{0,-1} + G_{1,0} H_{0,-1} - G_{1,\frac{1}{2+1}} H_{0,-1} - G_{1,-\frac{1+1}{2+1}} H_{0,-1} \\
 - & G_{-1,\frac{1}{2+1}} H_{0,-1} + G_{0,0} H_{0,-1} + G_{0,1} H_{0,-1} - G_{0,\frac{1+1}{2+1}} H_{0,-1} - G_{0,-\frac{1+1}{2+1}} H_{0,0} \\
 + & G_{1,-\frac{1+1}{2+1}} H_{0,0} - G_{1,0} H_{0,0} + G_{-1,\frac{1}{2+1}} H_{0,1} + G_{-1,-\frac{1+1}{2+1}} H_{0,1} - G_{0,\frac{1}{2+1}} H_{0,1} + G_{0,-\frac{1+1}{2+1}} H_{0,1} \\
 + & G_{1,0} H_{0,1} - G_{1,\frac{1}{2+1}} H_{0,1} - G_{1,-\frac{1+1}{2+1}} H_{0,1} - G_{-1,0} H_{1,-1} + G_{-1,\frac{1}{2+1}} H_{1,-1} - G_{-1,-\frac{1+1}{2+1}} H_{1,-1} \\
 + & G_{0,0} H_{1,-1} + 2G_{0,\frac{1}{2}} H_{1,-1} - 2G_{0,1} H_{1,-1} - G_{0,\frac{1+1}{2+1}} H_{1,-1} - G_{0,-\frac{1+1}{2+1}} H_{1,-1} \\
 - & G_{1,\frac{1}{2+1}} H_{1,-1} + G_{1,-\frac{1+1}{2+1}} H_{1,-1} + G_{-1,0} H_{1,0} - G_{-1,\frac{1}{2+1}} H_{1,0} - G_{-1,-\frac{1+1}{2+1}} H_{1,0} + G_{0,\frac{1}{2+1}} H_{1,0} \\
 - & G_{0,-\frac{1+1}{2+1}} H_{1,0} - G_{1,0} H_{1,0} + G_{1,\frac{1}{2+1}} H_{1,0} + G_{1,-\frac{1+1}{2+1}} H_{1,0} - G_{-1,0} H_{1,1} + G_{-1,-\frac{1+1}{2+1}} H_{1,1} \\
 - & G_{-1,-\frac{1+1}{2+1}} H_{1,1} + G_{0,0} H_{1,1} + 2G_{0,\frac{1}{2}} H_{1,1} - 2G_{0,1} H_{1,1} - G_{0,\frac{1+1}{2+1}} H_{1,1} - G_{0,-\frac{1+1}{2+1}} H_{1,1}
 \end{aligned}$$

1. Find a finite basis

$$I_5(s, t, u, m) = \sum_{j=1}^{\text{finite}} c_j f_j(s, t, u, m)$$

- determined by **singularities**

2. Landau bootstrap

apply enough analytic constraints to uniquely fix all c_j

3. Analytic regression

fit the c_j numerically

Where are the singularities?

$$I_G(p) = (n_{\text{int}} - 1)! \int_0^\infty \prod_{e \in E_{\text{int}}(G)} d\alpha_e \int \prod_{c \in \hat{C}(G)} d^d k_c \frac{1}{(\ell + i\varepsilon)^{n_{\text{int}}}} \delta\left(1 - \sum_{e \in E_{\text{int}}(G)} \alpha_e\right)$$

A necessary condition for a singularity is that the *integrand* is singular ($\ell=0$)

$$\ell = \sum_{e \in E_{\text{int}}(G)} \alpha_e (q_e^2 - m_e^2) = 0$$

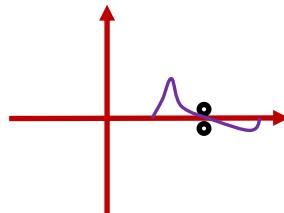
- every internal line is either on-shell ($q^2=m^2$) or $\alpha=0$ or both

A necessary condition for a singularity of the *integral* is that poles pinch the contour

for each loop k_c :

$$\sum_{e \in E_{\text{int}}(G^\kappa)} \alpha_e \frac{\partial}{\partial k_c} (q_e^2 - m_e^2) = 0.$$

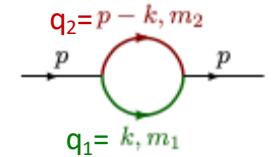
Double pole:



integration contour
pinched between poles

- since q_e are linear in k_c

$$\sum_{e \text{ in loop}} \pm \alpha_e q_e^\mu = 0$$



Landau loop equations

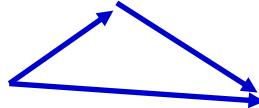
Coleman-Norton interpretation

Landau equations

$$\ell = \sum_{e \in E_{\text{int}}(G)} \alpha_e (q_e^2 - m_e^2) = 0$$

$$\sum_{e \text{ in loop}} \pm \alpha_e q_e^\mu = 0$$

4-momenta add up to zero after rescaling by α



[Coleman and Norton 1965]

Landau diagram is interpreted as space-time diagram

- momenta are on-shell (classical)
- α_e are the proper times for propagation

Gives some intuition: singularities due to classically allowed processes

Pham interpretation

- Pham 1967
- H. Hannesdottir, A. McLeod, MDS, C. Vergu [2211.07633](https://arxiv.org/abs/2211.07633)

Landau equations

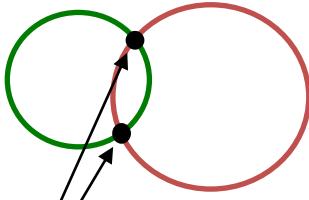
$$\ell = \sum_{e \in E_{\text{int}}(G)} \alpha_e (q_e^2 - m_e^2) = 0$$

$$\sum_{e \text{ in loop}} \pm \alpha_e q_e^\mu = 0$$

normal vectors
of on-shell constraints $q^2 = m^2$
are linearly dependent

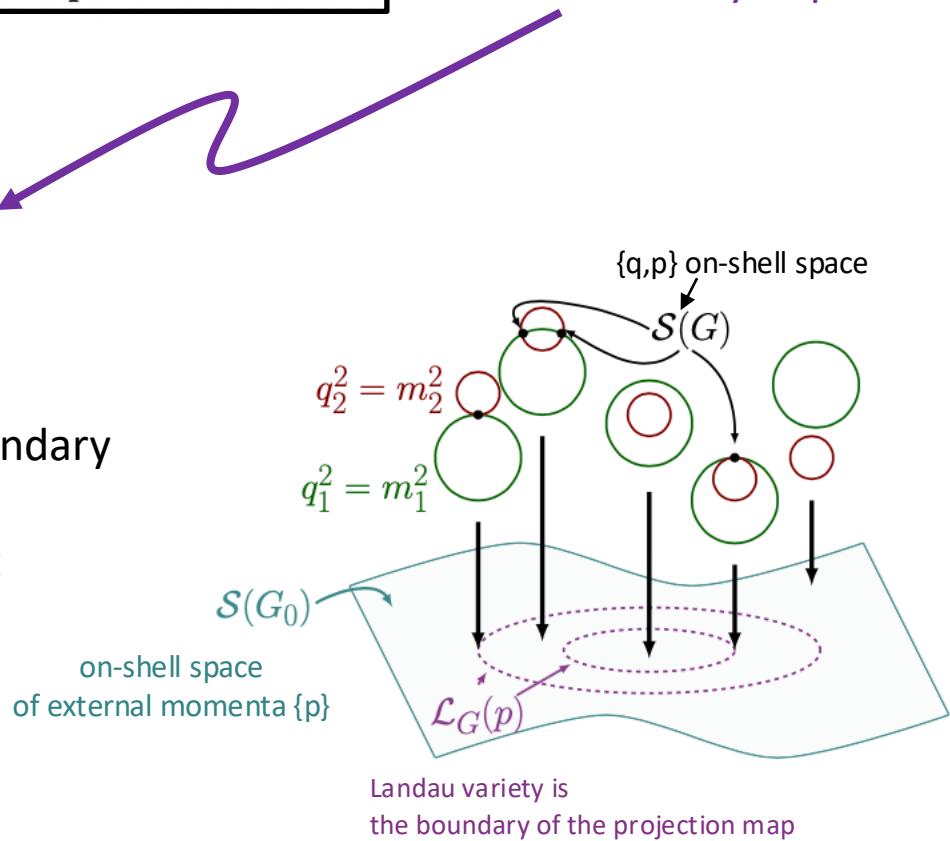
on-shell constraints (Euclidean $d=2$)

$$q_x^2 + q_y^2 = m_e^2$$



intersection
satisfies both
on-shell constraints

tangent on boundary
of space where
circles intersect



Bubble in d=2

Consider the simplest 1-loop diagram: the bubble in d=2

$$I_{\text{bubble}}(p) = \text{Diagram} = \int d^2k \frac{1}{k^2 - m^2 + i\varepsilon} \frac{1}{(p - k)^2 - m^2 + i\varepsilon} = \frac{-2\pi}{\sqrt{s(s - 4m^2)}} \ln \frac{\sqrt{4m^2 - s} - i\sqrt{s}}{\sqrt{4m^2 - s} + i\sqrt{s}}$$

The diagram shows a loop with a clockwise arrow. The incoming momentum is p and the outgoing momentum is p . The loop has a red label $p - k, m$ and a green label k, m .

- It has a **normal threshold** branch cut starting at $s=4m^2$
 - For $s > 4m^2$ the on-shell process $p \rightarrow p_1 + p_2$ is allowed for physical on-shell momenta

$$\text{Disc}[I(s)] = I(s + i\epsilon) - I(s - i\epsilon) = \frac{4\pi^2 i}{\sqrt{s(s - 4m^2)}}$$

Diagram: A horizontal line with an arrow labeled p enters from the left. It splits into two red arrows that diverge to the right.

- There is also a **pseudreshold** at $s=0$
 - There is a branch point at $s=0$ accessible with complex momenta
 - Does not correspond to anything physical happening

Bubble in d=2

More general 2D bubble: generic masses

$$I_{\text{O}}(p) = \text{---} \xrightarrow[p]{\text{---}} \text{---} = \lim_{\varepsilon \rightarrow 0^+} \int d^2 k \frac{1}{k^2 - m_1^2 + i\varepsilon} \frac{1}{(p - k)^2 - m_2^2 + i\varepsilon},$$

Going to Feynman parameters

$$I_{\text{O}}(s) = \lim_{\varepsilon \rightarrow 0^+} \int_0^1 d\alpha \frac{-i\pi}{s\alpha(1-\alpha) - m_1^2\alpha - m_2^2(1-\alpha) + i\varepsilon} = \ell$$

integrand is singular ($\ell = 0$) at

$$\alpha_{\pm} = \frac{s + m_2^2 - m_1^2 \pm \sqrt{s^2 - 2s(m_1^2 + m_2^2) + (m_1^2 - m_2^2)^2 + i s \varepsilon}}{2s}.$$

on-shell locus

- necessary but not sufficient condition for singularities of integral

singularities require pinches, i.e. $\frac{d\ell}{d\alpha} = 0$

normal threshold $s = (m_1 + m_2)^2 - i\varepsilon$,

pseudothreshold $s = (m_1 - m_2)^2 + i\varepsilon$,

two solutions

$$\left. \begin{aligned} \alpha_{\pm} &= \frac{m_2}{m_2 + m_1} + i\varepsilon \operatorname{sgn}(m_2 - m_1), \\ \alpha_{\pm} &= \frac{m_2}{m_2 - m_1} - i\varepsilon \operatorname{sgn}(m_2 - m_1). \end{aligned} \right\}$$

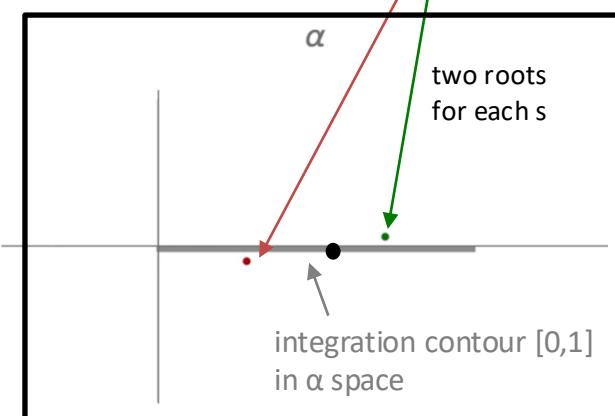
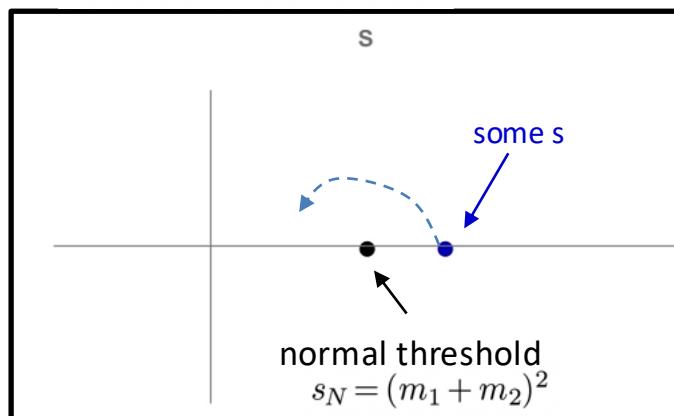
- location of branch points
- solutions to Landau equations

Homotopy and Homology

$$I_{\text{O}}(s) = \lim_{\varepsilon \rightarrow 0^+} \int_0^1 d\alpha \frac{-i\pi}{s\alpha(1-\alpha) - m_1^2\alpha - m_2^2(1-\alpha) + i\varepsilon} \\ = \int_0^1 \frac{d\alpha}{[\alpha - \alpha_+(s)][\alpha - \alpha_-(s)]}$$

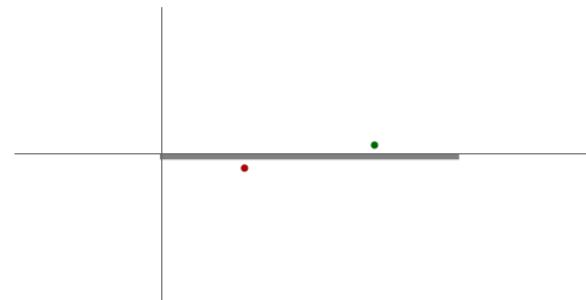
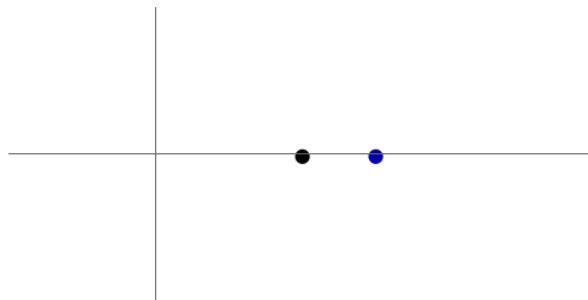
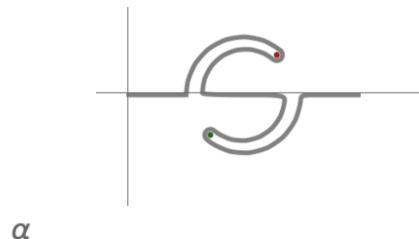
on-shell locus: $\alpha = \alpha_{\pm}$

$$\alpha_{\pm} = \frac{s + m_2^2 - m_1^2 \pm \sqrt{s^2 - 2s(m_1^2 + m_2^2) + (m_1^2 - m_2^2)^2 + i\varepsilon}}{2s}.$$



What happens as we take a monodromy (discontinuity) of s around s_N ?

- Poles α_{\pm} move around too
- Contour must move out of the way to avoid poles
- After full loop, can use Leray residue theorem to get discontinuity



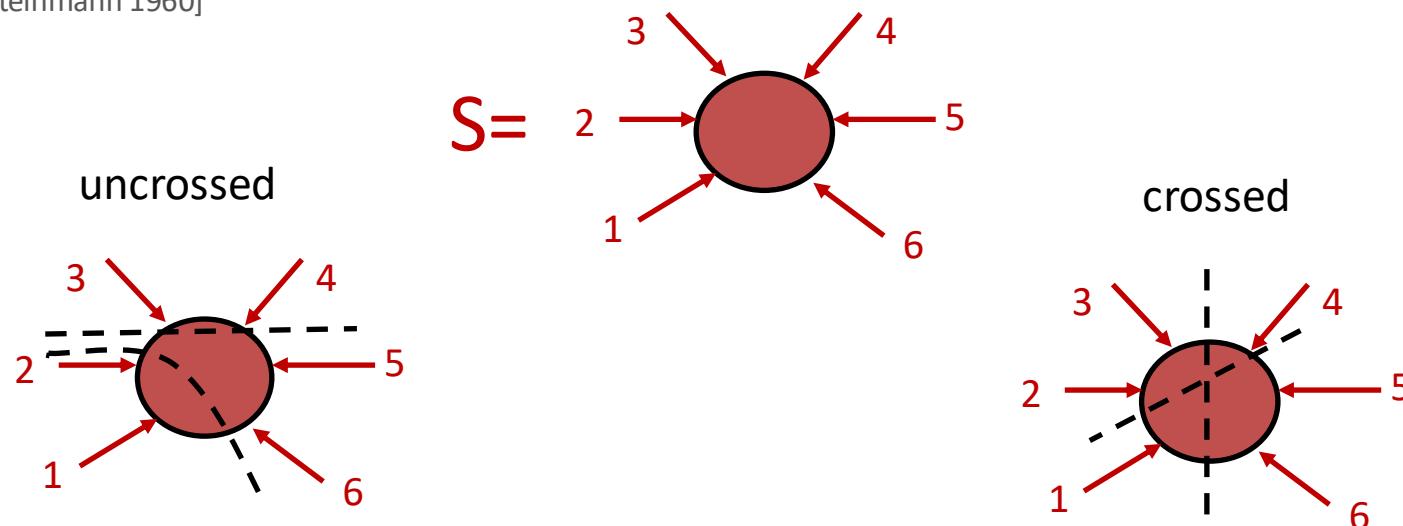
Sequential discontinuities

- Can take a discontinuity of a discontinuity

$$\text{Disc}_N \text{Disc}_P I = \frac{-8\pi^2 i}{\sqrt{(s - s_N)(s - s_P)}}$$

Steinman relations : S matrix cannot have sequential discontinuities in crossed channels

[Steinmann 1960]



possible term: $\ln(p_1 + p_2)^2 \ln(p_3 + p_4)^2$

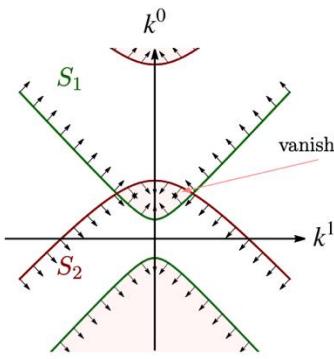
cannot appear $\ln(p_1 + p_2 + p_3)^2 \ln(p_2 + p_3 + p_4)^2$

- Algebraic geometry need to understand contour deformations and pinches
- Gets pretty complicated...

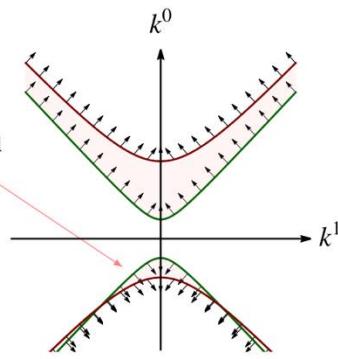
Different kinds of singularities

Simple pinches

near normal threshold



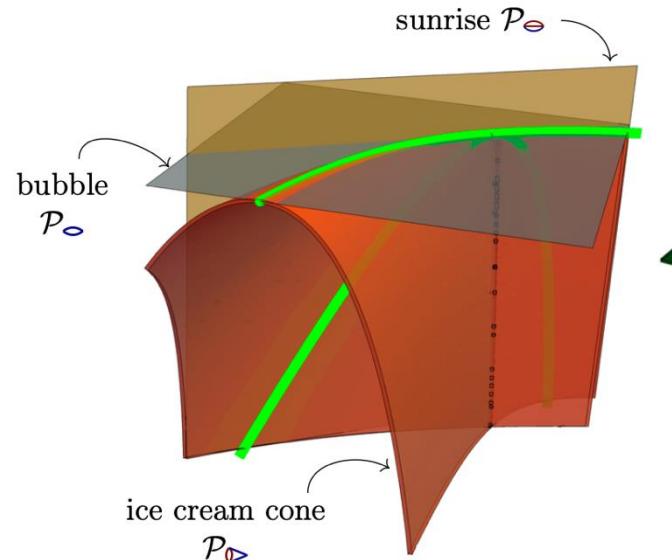
near pseudonormal threshold



- hypersurfaces meet transversely
- e.g. physical thresholds

McLeod, Hanesdottir, MDS, Vergu
arXiv:2211.07633

Non-simple pinches



- hypersurfaces meet tangentially
 - e.g. sunrise in the ice-cream cone

$$G_0 = \begin{array}{c} \diagup \diagdown \\ q_4 \end{array} \begin{array}{c} \diagup \diagdown \\ q_3 \end{array} \begin{array}{c} \diagup \diagdown \\ q_2 \end{array} \xrightarrow{\bar{\kappa}} G_0 = \begin{array}{c} \diagup \diagdown \\ q_1 \end{array} \begin{array}{c} \diagup \diagdown \\ q_2 \end{array}$$

- Permanent pinches (e.g. IR divergences)
- Pinches at infinity

Symbol

The symbol is an efficient way to represent polylogarithmic Feynman integrals

[Goncharov, Spradlin, Vergu, Volovich PRL 2010]

- Used symbol to simplify 17 page 2-loop 6 point amplitude to a few lines



15th anniversary!

etc.

Symbol is a map that extracts the dlog forms

$$\mathcal{S} \left[\int_a^b d \ln R_1 \circ \cdots \circ d \ln R_n \right] = R_1 \otimes \cdots \otimes R_n$$

$$\mathcal{S} [\ln x \ln y] = x \otimes y + y \otimes x$$

$$\mathcal{S} [\text{Li}_n(x)] = -(1-x) \otimes \underbrace{x \otimes \cdots \otimes x}_{n-1}$$

$$\mathcal{S} \left[\frac{1}{n!} \ln^n x \right] = \underbrace{x \otimes \cdots \otimes x}_n$$

Advantages

- Efficient way to represent integrals
- Possible logarithmic singularities \Leftrightarrow symbol “letters” vanish
- Transparently encodes discontinuities and sequential discontinuities

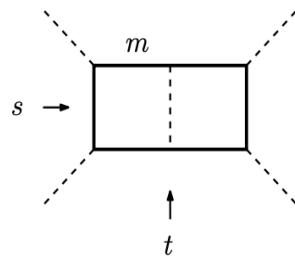
Constructing a Finite basis

The **symbol** is the key to making the the problem finite

1. Identify singularities

- All solution to Landau equations
- Some singularities may be on higher Riemann sheets

outer mass double box



physical singularities

$$s = 4m^2, \quad s \rightarrow \infty,$$

$$t = 4m^2, \quad t \rightarrow \infty,$$

$$m^2 = 0,$$

$$u = -\frac{4m^2}{s}, \quad v = -\frac{4m^2}{t}.$$

unphysical singularities

$$s = 0, \quad t = 0, \quad s + t = 0,$$

$$st + 4m^2s + 4m^2t = 0.$$

2. Identify possible letters

- algebraic functions (with square roots) of singularities with no new singularities

$$\beta_u = \sqrt{1+u}, \quad \beta_v = \sqrt{1+v}, \quad \beta_{uv} = \sqrt{1+u+v}.$$

$$\begin{aligned} L_1 &= u, & L_2 &= v, & L_7 &= \frac{\beta_v - 1}{\beta_v + 1}, & L_8 &= \frac{\beta_{uv} - 1}{\beta_{uv} + 1}, \\ L_3 &= 1 + u, & L_4 &= 1 + v, & L_9 &= \frac{\beta_{uv} - \beta_u}{\beta_{uv} + \beta_u}, & L_{10} &= \frac{\beta_{uv} - \beta_v}{\beta_{uv} + \beta_v}, \\ L_5 &= u + v, & L_6 &= \frac{\beta_u - 1}{\beta_u + 1}, & L_{11} &= 1 + u + v, \end{aligned}$$

3. Finite basis

- sum over all possible symbols using the alphabet

$$\mathcal{S}[\mathcal{I}] = \sum c_{\{i\}} L_{i_1} \otimes L_{i_2} \otimes L_{i_3} \otimes L_{i_4}$$

Landau Bootstrap

A. Construct finite basis

1. Find singularities
2. Find letters
3. Construct finite basis

MDS, Hannesdottir, McLeod, Vergu [2410.02424](#)

$$\mathcal{S}[\mathcal{I}] = \sum c_{\{i\}} L_{i_1} \otimes L_{i_2} \otimes L_{i_3} \otimes L_{i_4}$$

B. Fix the coefficients

- Integrability
- Galois symmetry
- α -positivity
- First-entry conditions
- Last-entry conditions
- Geneological constraints
 - Sequential discontinuities
 - Cluster adjacency conditions
- Regions analysis
- Direct calculation

C. Fix the non-symbol rational prefactor

Integrability

- Every interated dlog integral has a symbol

$$f(u, v; \Gamma) = c_{i_1, i_2, \dots, i_n} \int_{\Gamma} d \ln L_{i_1} \circ \dots \circ d \ln L_{i_n} \cdot \xrightarrow{\quad} \sum c_{i_1, i_2, \dots, i_n} L_{i_1} \otimes \dots \otimes L_{i_n}$$

• Not every symbol corresponds to a function

Learning the Simplicity of Scattering Amplitudes
 Clifford Cheung (Caltech), Aurélien Dersy (Harvard U. and IAFI, Cambridge),
 Matthew D. Schwartz (Harvard U. and IAFI, Cambridge) (Aug 8, 2024)
 e-Print: [2408.04720](https://arxiv.org/abs/2408.04720) [hep-th]

- For $f(u, v, \Gamma)$ to be a function, must be independent of local path deformations

$$\xrightarrow{\quad} [\partial_u, \partial_v] f = 0$$

- Derivatives only act on the last entry of the symbol (end of integration contour)

$$\partial_u [S \otimes K \otimes L] = (\partial_u \ln L) [S \otimes K]$$

$$\partial_v \partial_u [S \otimes K \otimes L] = (\partial_v \partial_u \ln L) [S \otimes K] + (\partial_u \ln L) (\partial_v \ln K) S$$

$$[\partial_u, \partial_v] [S \otimes K \otimes L] = \underbrace{[(\partial_u \ln L) (\partial_v \ln K) - (\partial_v \ln K) (\partial_u \ln L)] S}_{\text{must vanish}} \quad (\text{integrability condition})$$

α positivity

- Symbol encodes all branch points, even on unphysical sheets

$$I_G(p) = (n_{\text{int}} - 1)! \int_0^\infty \prod_{e \in E_{\text{int}}(G)} d\alpha_e \int \prod_{c \in \hat{C}(G)} d^d k_c \frac{1}{(\ell + i\varepsilon)^{n_{\text{int}}}} \delta \left(1 - \sum_{e \in E_{\text{int}}(G)} \alpha_e \right)$$

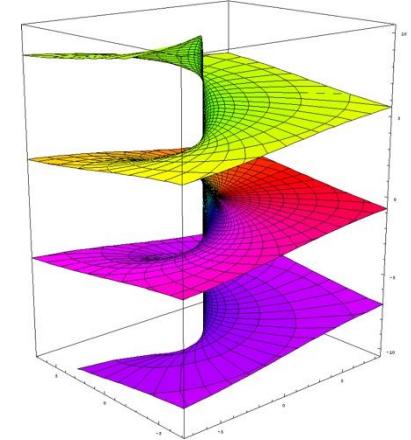


actual Feynman integral on physical sheet is over positive α

- Discontinuities/monodromies act on first entry of the symbol

$$I = \int \omega_1 \int \omega_2 \dots \int \omega_n$$

$$dI = \omega_1 \int \omega_2 \dots \int \omega_n$$



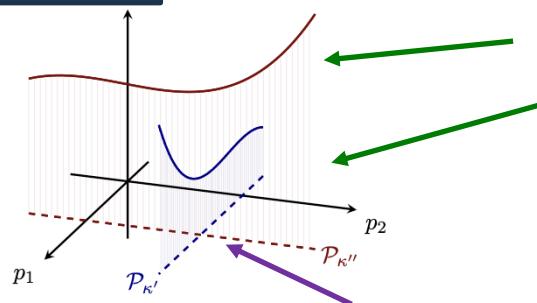
- Singularity for physical momenta (physical sheet) \leftrightarrow singularity of first entry
 - See if $\alpha > 0$ in solutions to Landau equations: constrain first symbol entries

Genealogical constraints

Which symbol entries can be adjacent?

- Adjacent means sequential discontinuities are possible

Pham approach

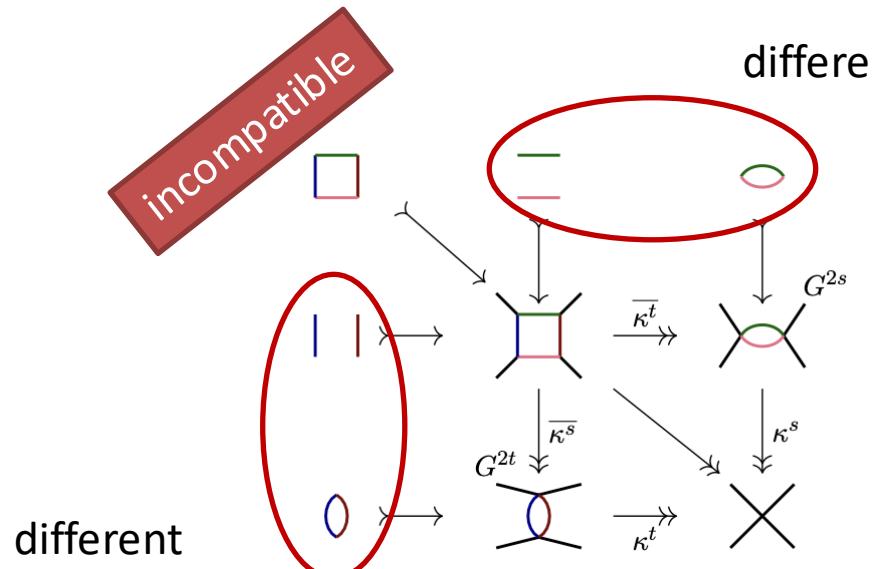
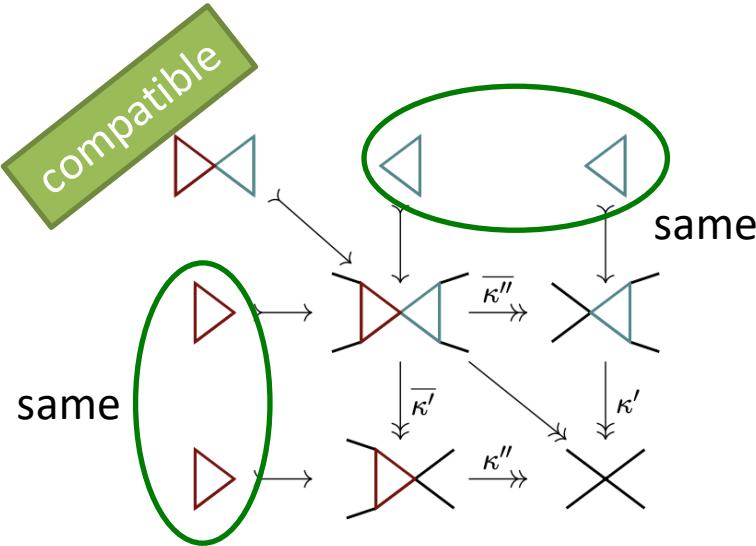


on-shell surfaces may not intersect in *internal momenta*

- vanishing cell from first monodromy doesn't intersect integration contour of second
- sequential monodromy vanishes

$$(\mathbb{1} - \mathcal{M}_{P_{\kappa'}})(\mathbb{1} - \mathcal{M}_{P_{\kappa''}})I_G(p) = 0$$

Singular surfaces intersect transversally in *external momenta*

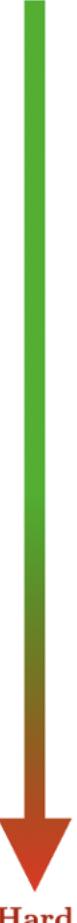


- e.g. Steinmann relations

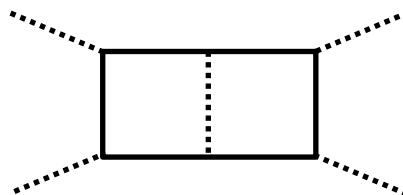
Outer-mass double box

Landau bootstrap

Easy



Constraints	# Coeffs
All Symbols	20736
Integrability	6993
Galois symmetry	861
Physical branch cuts	161
Genealogical constraints	28
α -positive thresholds	6



Caron-Huot & Henn '14

- Weight 4 symbol, 12 Letters
- Work from easy generic stuff (integrability) to hard

$$\begin{aligned}
 S(\tilde{I}_{\text{dbox}}) = & -L_6 \otimes \frac{L_1}{L_3} \otimes L_6 \otimes L_9 - L_6 \otimes \frac{L_1}{L_3} \otimes L_9 \otimes L_6 \\
 & + L_6 \otimes L_6 \otimes \frac{L_1 L_2}{L_3 L_5} \otimes L_9 + L_6 \otimes L_9 \otimes \frac{L_2}{L_5} \otimes L_6 \\
 & + L_6 \otimes L_6 \otimes L_8 \otimes L_6 + L_6 \otimes L_9 \otimes L_8 \otimes L_9 \\
 & + L_7 \otimes L_{10} \otimes \frac{L_2}{L_5} \otimes L_6 + L_7 \otimes L_{10} \otimes L_8 \otimes L_9 \\
 & + L_7 \otimes L_7 \otimes \frac{L_1}{L_5} \otimes L_9 + L_7 \otimes L_7 \otimes L_8 \otimes L_6 .
 \end{aligned}$$



- Agrees with Caron-Huot & Henn '14
- Also need **rational prefactor** -- not in the symbol
 - Compute from maximal cut



Landau bootstrap works!

Top down or bottom up?

1. Top down:
Landau bootstrap



Easy

Hard

Constraints **# Coeffs**

All Symbols	20736
Integrability	6993
Galois symmetry	861
Physical branch cuts	161
Genealogical constraints	28
α -positive thresholds	6

Hard

Easy



2. Bottom up:
Analytic regression with
lattice reduction

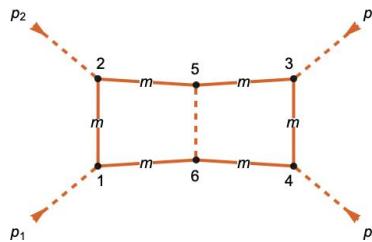
Solving the Landau equations

Lots of ways to solve the Landau equations

- Solve them by hand (e.g. Eden et al 1950)
- HyperInt (Panzer 2014)
- PLD (Fevola, Mizera, Telen 2013)
- BaikovLetter (Jiang et al 2024)
- Recursive approach (Caron-Huot, Correia and Giroux 2024)
- Numerical implementation for any diagram (Correia, Giroux, Mizera 2024) **SOF~~EA~~A**

Input

```
diag =  
  {{{{1, 2}, m1}, {{2, 5}, m2}, {{3, 5}, m3}, {{3, 4}, m4}, {{4, 6}, m5},  
   {{1, 6}, m6}, {{5, 6}, m7}}, {{1, M1}, {2, M2}, {3, M3}, {4, M4}}}//.  
 m7 → 0 // . m_ → m // . M_ → 0  
FeynmanPlot[diag]  
  
{{{1, 2}, m}, {{2, 5}, m}, {{3, 5}, m}, {{3, 4}, m}, {{4, 6}, m},  
 {{1, 6}, m}, {{5, 6}, 0}}, {{1, 0}, {2, 0}, {3, 0}, {4, 0}}}
```

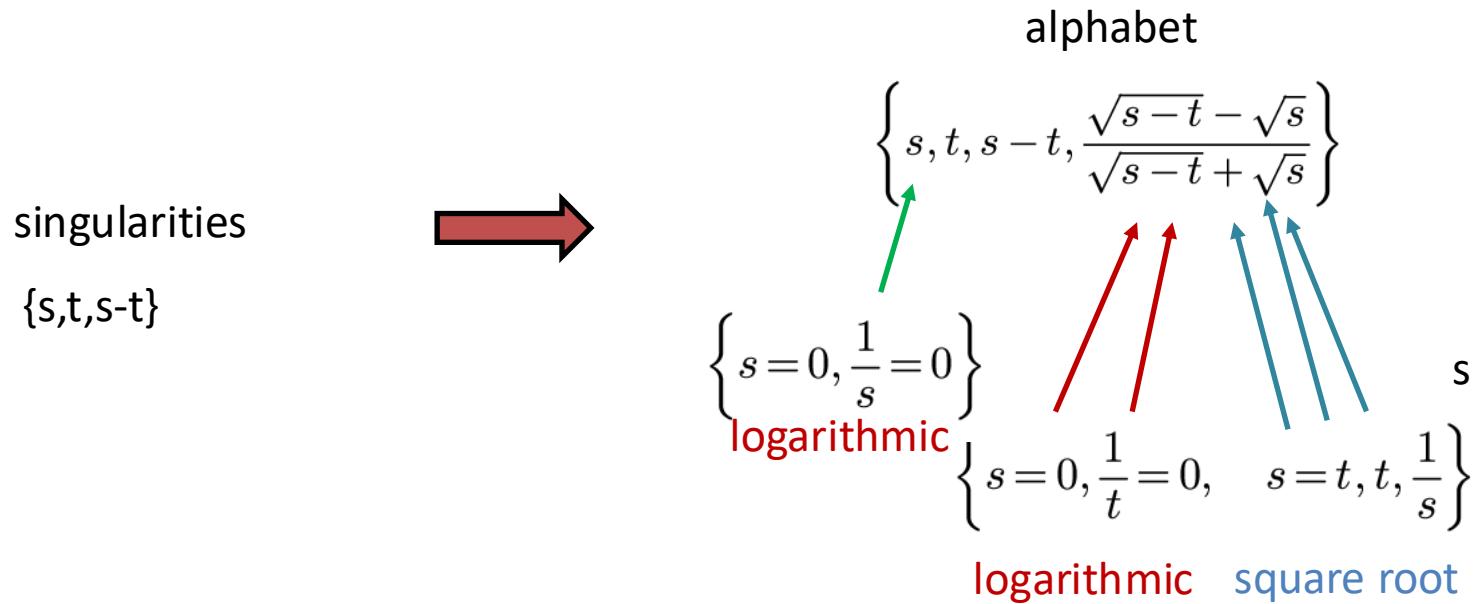


Output

```
]:= candidateSingularities =  
  SOFIA[diag, SolverBound → Infinity];  
 % // TableForm  
  
⌚ Total runtime= 2.86441  
  
:=  
 mm  
 s12  
 s23  
 s12 + s23  
 mm - s12  
 mm - s23  
 4 mm - s12  
 4 mm - s23  
 4 mm s12 + 4 mm s23 - s12 s23  
 mm2 s12 - 2 mm s12 s23 - 4 mm s232 + s12 s232  
 mm s122 + 4 mm s12 s23 - s122 s23 + 4 mm s232 - s12 s232
```

Finds all
singularities
(simple & non-simple)

Singularities to letters



Landau's original paper determined if physical singularities were logarithmic or square-root

- In general, singularities may appear multiple times on multiple sheets
- Still true that singularities are always either **logarithmic** or **square root**
- Automated codes like SOFIA seem to be able to tell logarithmic from square root

Finite basis

Location of Landau singularities

$$\begin{aligned}
 s &= 4m^2, \quad s \rightarrow \infty, \\
 t &= 4m^2, \quad t \rightarrow \infty, \\
 m^2 &= 0, \\
 s &= 0, \quad t = 0, \quad m^2 \rightarrow \infty, \\
 s + t &= 0, \quad st - 4m^2s - 4m^2t = 0.
 \end{aligned}$$

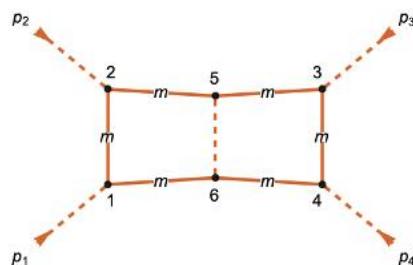


symbol alphabet

$$\tilde{A} = \left\{ u, v, 1+u, 1+v, u+v, 1+u+v, \frac{\beta_u - 1}{\beta_u + 1}, \frac{\beta_v - 1}{\beta_v + 1}, \frac{\beta_{uv} - 1}{\beta_{uv} + 1}, \frac{\beta_{uv} - \beta_u}{\beta_{uv} + \beta_u}, \frac{\beta_{uv} - \beta_v}{\beta_{uv} + \beta_v}, \frac{\beta_{uv} - \beta_u \beta_v}{\beta_{uv} + \beta_u \beta_v} \right\}$$

$$\beta_u = \sqrt{1+u}, \beta_v = \sqrt{1+v}, \beta_{uv} = \sqrt{1+u+v}.$$

- SOFIA can also produce the alphabet (may be larger than needed)
- Symbol weight $\leq 2L$, L = loop order
 - 12 letters
 - $12^4 = 20,736$ symbol entries

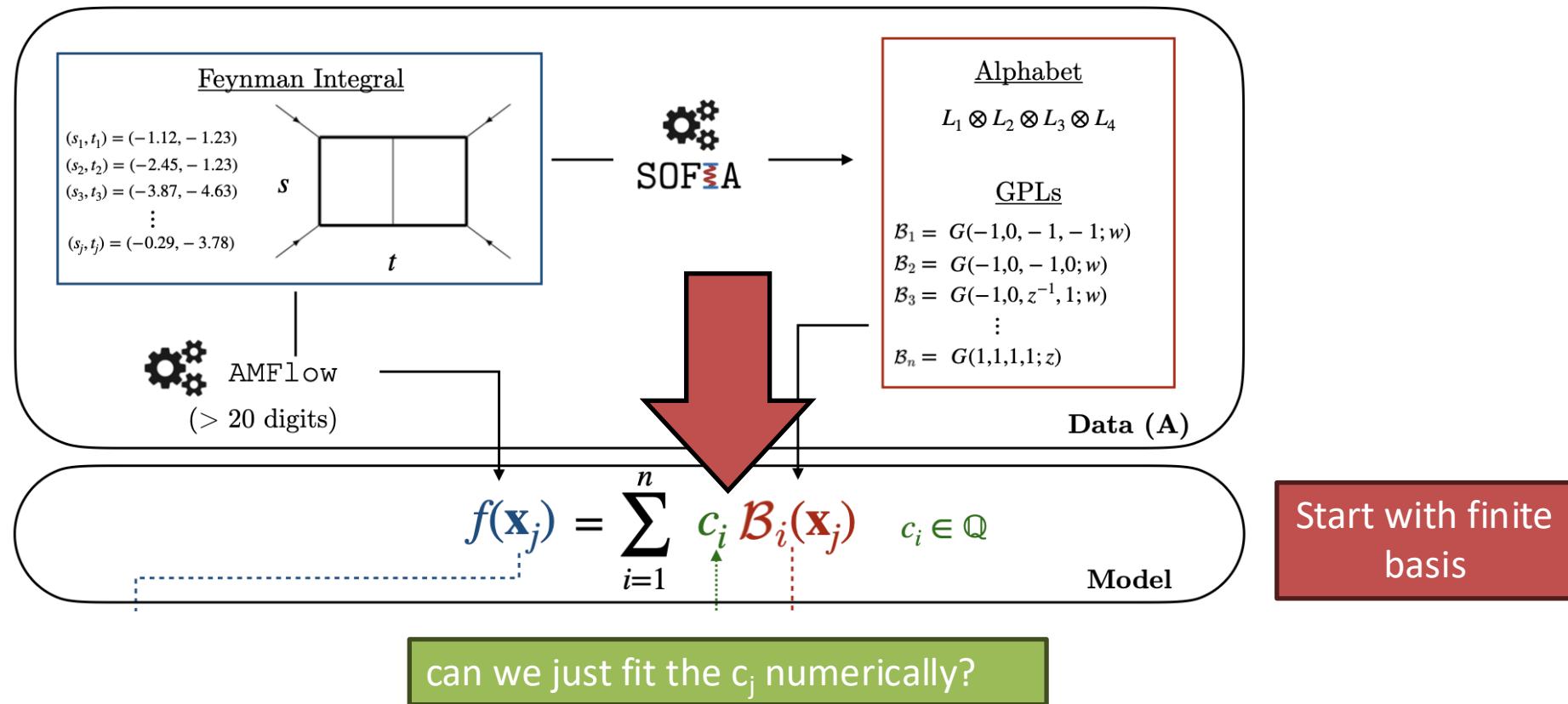


2 loop outer-mass double box

$$\mathcal{S}[\mathcal{I}] = \sum c_{\{i\}} L_{i_1} \otimes L_{i_2} \otimes L_{i_3} \otimes L_{i_4}$$

Finite basis!

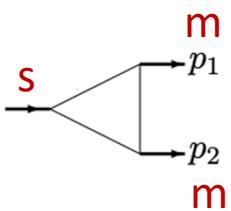
Numerical bootstrap



Requirements

1. Evaluate $f(x)$ numerically to high precision
2. Evaluate $B(x)$ to high precision
3. Solve the equation without losing precision

1-loop triangle

exact answer: $T =$  $= \frac{x}{m^2(1-x^2)} \left(-\frac{\pi^2}{3} + 4G(0, -1, x) - 2G(0, 0; x) \right)$

Alphabet (from Sofia)

one invariant $x = \frac{1 - \sqrt{1 + \frac{4m^2}{s}}}{1 + \sqrt{1 + \frac{4m^2}{s}}}$

3 letters

$$A = \{L\} = \{m, s, 4m^2 - s\} \leftrightarrow \{x, 1 + x, 1 - x\}$$

Basis is

$$T(x) = \sum_{i=1}^{17} c_i \mathcal{B}_i(x) = P(x) \times [c_{ij} L_i \otimes L_j + (a_j + b_j \pi) L_j + d_1 \pi^2 + d_2 \zeta_3]$$

rational prefactor (Sofia computes)
Total runtime = 0.042703

$\frac{1}{16 \pi^4 \sqrt{(4m^2 - s) s}}$

9 terms weight 2 6 terms weight 1 2 numbers

17 constants

- Need $T(x)$ and $\mathcal{B}_i(x)$ to high precision

Numerical Feynman Integration

- AMFlow: numerical evaluation for Feynman loop integrals to high precision

Integration-by-Parts (IBP) [Chetyrkin, Tkachov, 1981]

$$\mathcal{I}(p_i, m_e) = \int \prod_{j=1}^L d^D k_j \frac{N(p_i, k_j)}{\prod_{e=1}^E (q_e^2 - m_e^2 + i\epsilon)} \Rightarrow \mathcal{I} = \sum_a R_a(p_i \cdot p_j, m_e) \text{MI}_a$$

Rational functions Master integrals
(basis)

Auxiliary mass flow method (AMFlow)

[Liu, Ma, 1801.10523; 2107.01864;
2201.11669; 2201.11637]

1. Introduce an auxiliary mass η to some of the propagator denominators
2. Set up closed differential equations w.r.t η using IBPs
3. Solve the differential equations numerically with boundary conditions $\eta \rightarrow \infty$

Timing on 1-loop triangle

- each phase space point takes 5-10 CPU-min for 30 significant digits

Numerical basis functions

$$T(x) = \sum_{i=1}^{17} c_i \mathcal{B}_i(x) = P(x) \times [c_{ij} L_i \otimes L_j + (a_j + b_j \pi) L_j + d_1 \pi^2 + d_2 \zeta_3]$$

- Rational prefactor evaluation is instant
- Weight-2 symbols can be integrated into closed form expressions
 - FiberSymbol in Polylogtools can do this
 - Numerical evaluation is very fast

Integrals of higher weight symbols are not always known

- Can always integrate numerically along a path

$$\tilde{B} = \sum_{i=1}^{|\tilde{A}|} c_{i_1, i_2, i_3, i_4} \int_0^1 d \log L_{i_4}(\lambda_4) \int_0^{\lambda_4} d \log L_{i_3}(\lambda_3) \int_0^{\lambda_3} d \log L_{i_2}(\lambda_2) \int_0^{\lambda_2} d \log L_{i_1}(\lambda_1),$$

- Individual terms may be path dependent, but final result is integrable
- Can do the first and last integral analytically
 - Reduce weight by 2: speeds up integration tremendously
- Sometimes analytic integrals can be so complicated that it is faster to do the integral numerically
 - More work needed on efficient numerical evaluation

Matrix inversion

- Evaluate both sides of this equation at 17 points
- Solve 17 linear equations for coefficients
 - i.e. invert the matrix

$$T(x) = \sum_{i=1}^{17} c_i \mathcal{B}_i(x)$$

$$M_{ij} = \mathcal{B}_i(\mathbf{x}_j) \quad \rightarrow \quad c_i = (M^{-1})_{ij} \cdot T(x_j)$$

Problems:

1. No way to impose that c_j are rational
 - Cannot ever resolve functions that differ by irrational constants

$$2[\ln x] + 3[\pi \ln x] = (1 + \pi)[\ln x] + \left(2 + \frac{1}{\pi}\right)[\pi \ln x]$$

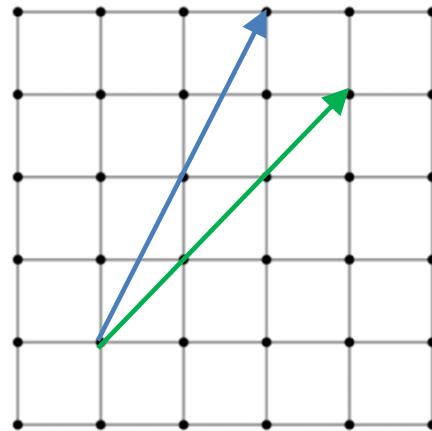
2. Matrix inversion loses precision very fast
 - Controlled by condition number

$$\kappa(M) = \|M\| \|M^{-1}\|.$$

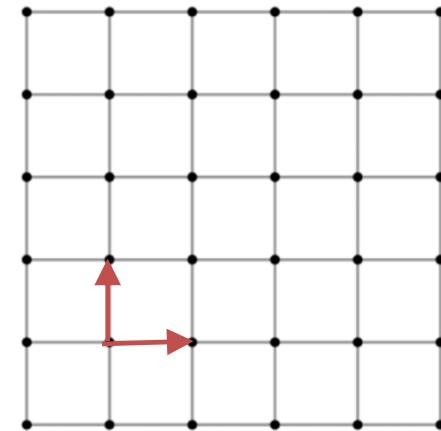
- Generally our matrices will be ill-conditioned
 - (they come from smooth functions)

Lattice reduction

vectors span
a lattice



$$\vec{u}_1^2 + \vec{u}_2^2 = 24$$



$$\vec{v}_1^2 + \vec{v}_2^2 = 2$$

- Lattice reduction finds another set of vectors for same lattice
- Can minimize some norm (length of lattice vectors)
- NP-Hard problem: no polynomial-time algorithm for truly best solution
- Efficient algorithms exist to find what is almost always the minimum

Lattice reduction

Rational number coefficients can be fit for numbers using lattice reduction

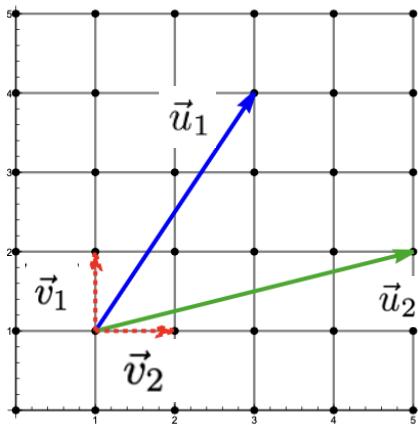
$$f = - \int_0^1 du \int_0^1 dv \frac{\log(1 - uv) + v \log(1 - u)}{uv} = \frac{\pi^2}{6} + \zeta_3 = c_1 \pi^2 + c_2 \zeta_3 = \vec{c} \cdot \vec{v}$$

$$f=2.847 \quad \longleftrightarrow \quad \text{Assume 4 digits known} \quad \longleftrightarrow \quad \pi^2=9.870 \quad \zeta_3=1.202$$

Multiply by 10^3 and put into a matrix

$$\begin{pmatrix} \vec{u}_1 \\ \vec{u}_2 \\ \vec{u}_3 \end{pmatrix} = \begin{pmatrix} 10^3 f & 1 & 0 & 0 \\ 10^3 \pi^2 & 1 & 0 & 0 \\ 10^3 \zeta_3 & 1 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 2847 & 1 & 0 & 0 \\ 9870 & 0 & 1 & 0 \\ 1202 & 0 & 0 & 1 \end{pmatrix} \xrightarrow{\text{lattice reduction}} \begin{pmatrix} \vec{v}_1 \\ \vec{v}_2 \\ \vec{v}_3 \end{pmatrix} = \begin{pmatrix} 0 & -6 & 1 & 6 \\ 3 & -11 & 5 & -15 \\ 62 & 4 & -2 & 7 \end{pmatrix}$$

vectors span a lattice



- Lattices are the same so v in the span of u

$$\vec{v}_1 = -6\vec{u}_1 + \vec{u}_2 + 6\vec{u}_3.$$

First component

$$0 = 10^3 \times (-6f + \pi^2 + 6\zeta^3) \quad \checkmark$$

Precision requirements

Rational number coefficients can be fit for numbers lattice reduction

$$f = - \int_0^1 du \int_0^1 dv \frac{\log(1 - uv) + v \log(1 - u)}{uv} = \frac{\pi^2}{6} + \zeta_3 = c_1 \pi^2 + c_2 \zeta_3 = \vec{c} \cdot \vec{v}$$

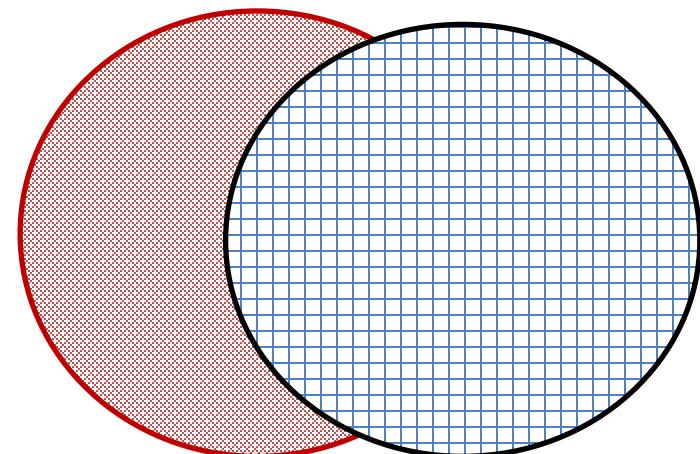
Q: how many digits of f are required to fit rational c_1 and c_2 ?

$$\vec{v} = (\pi^2, \zeta_3)$$

- Multiple solutions implies that $\vec{c} \cdot \vec{v} = 0$
- Multiply by GCD so c_1 and c_2 are integers
 - There are $(10^R)^n = 10^{2R}$ vectors (c_1, c_2)
- Assume d digits of precision on v
 - $\vec{c} \cdot \vec{v}$ produce $10^{2R} d$ digit numbers
- There are only 10^d d -digit numbers in all
- Need precision $d > nR$ to fit pure numbers
 - Information content must be sufficient

assume size of c 's

$$c_1 \sim c_2 \lesssim 10^R$$



Precision requirements

For fitting *functions* we can sample at multiple points

$$f(x) = G(0, 1; x) - G(1, -1; x) \quad x_1 = 4/10, x_2 = 9/10$$

$$\mathcal{B}(x) = \{G(1, 0; x), G(0, 1; x), G(0, -1; x), G(1, -1; x)\}$$

$$M = \text{round } 10^s \left(\begin{array}{c|c} f(\mathbf{x}_1) \cdots f(\mathbf{x}_p) & \\ \mathcal{B}_1(\mathbf{x}_1) \cdots \mathcal{B}_1(\mathbf{x}_p) & \\ \vdots & \\ \mathcal{B}_n(\mathbf{x}_1) \cdots \mathcal{B}_n(\mathbf{x}_p) & \end{array} \middle| 10^{-s} \mathbb{I}_{n+1} \right) = \begin{pmatrix} \mathbf{v}_1 \\ \mathbf{v}_2 \\ \mathbf{v}_3 \\ \mathbf{v}_4 \\ \mathbf{v}_5 \end{pmatrix} = \begin{pmatrix} -35 & -24 & 1 & 0 & 0 & 0 & 0 \\ 92 & 154 & 0 & 1 & 0 & 0 & 0 \\ -45 & -129 & 0 & 0 & 1 & 0 & 0 \\ 36 & 75 & 0 & 0 & 0 & 1 & 0 \\ -10 & -106 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

- With p points and d digits
 - net digits of information is $p \times d$
- Expected digits needed

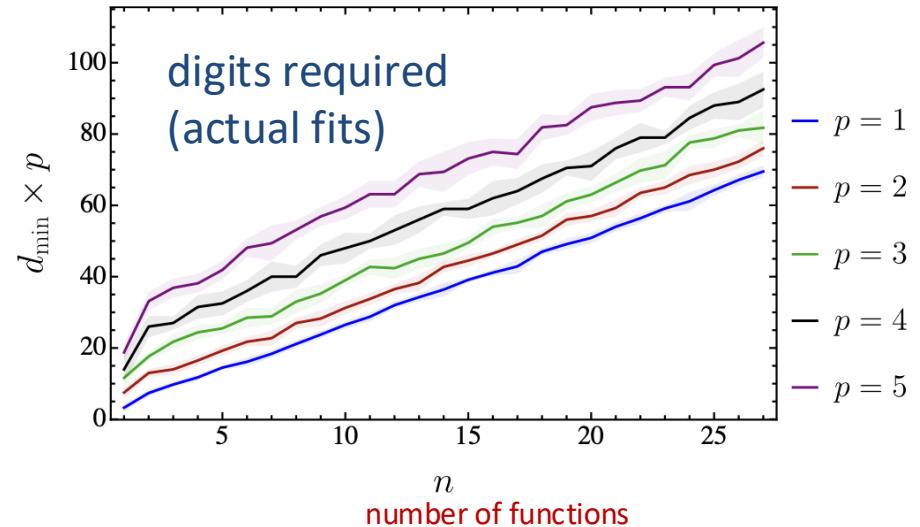
$$d \gtrsim \frac{nR}{p}$$

digits of precision req'd

basis functions

size of integers

number of points

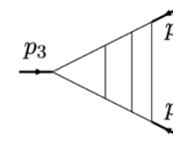
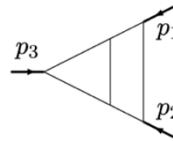
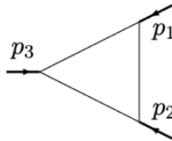


We can trade off digits of precision for number of points

- Will work even if only a few digits of precision are available!

Example 1: Triangles

Triangle ladder diagrams



exact results known

$$T_1(z) = \frac{1}{z - \bar{z}} \left[2\text{Li}_2(z) - \text{Li}_2(\bar{z}) + \log(z\bar{z}) \log \left(\frac{1-z}{1-\bar{z}} \right) \right],$$

$$T_2(z) = \frac{1}{(1-z)(1-\bar{z})(z-\bar{z})} \left[6\text{Li}_4(z) - 6\text{Li}_4(\bar{z}) - 3 \log(z\bar{z}) (\text{Li}_3(z) - \text{Li}_3(\bar{z})) + \frac{1}{2} \log^2(z\bar{z}) (\text{Li}_2(z) - \text{Li}_2(\bar{z})) \right],$$

$$T_3(z) = \frac{1}{(1-z)^2(1-\bar{z})^2(z-\bar{z})} \left[20\text{Li}_6(z) - 20\text{Li}_6(\bar{z}) - 10 \log(z\bar{z}) (\text{Li}_5(z) - \text{Li}_5(\bar{z})) + \log^2(z\bar{z}) (\text{Li}_4(z) - \text{Li}_4(\bar{z})) - \frac{1}{6} \log^3(z\bar{z}) (\text{Li}_3(z) - \text{Li}_3(\bar{z})) \right]$$

Allow all possibilities up to weight 6

Weight-0: 1

Weight-1: $G(a_1, x)$, π

Weight-2: $G(a_1, a_2, x)$, $\pi \times G(a_1, x)$, ζ_2

Weight-3: $G(a_1, a_2, a_3, x)$, $\pi \times G(a_1, a_2, x)$, $\zeta_2 \times G(a_1, x)$, ζ_3

Weight-4: $G(a_1, a_2, a_3, a_4, x)$, $\pi \times G(a_1, a_2, a_3, x)$, $\zeta_2 \times G(a_1, a_2, x)$, $\pi^3 \times G(a_1, x)$, $\zeta_3 \times G(a_1, x)$, ζ_4

Weight-5: $G(a_1, a_2, a_3, a_4, a_5, x)$, $\pi \times G(a_1, a_2, a_3, a_4, x)$, $\zeta_2 \times G(a_1, a_2, a_3, x)$, $\pi^3 \times G(a_1, a_2, x)$, $\zeta_3 \times G(a_1, a_2, x)$, $\zeta_4 \times G(a_1, x)$, ζ_5 , $\zeta_2 \times \zeta_3$

Weight-6: $G(a_1, a_2, a_3, a_4, a_5, a_6, x)$, $\pi \times G(a_1, a_2, a_3, a_4, a_5, x)$, $\zeta_2 \times G(a_1, a_2, a_3, a_4, x)$, $\pi^3 \times G(a_1, a_2, a_3, x)$, $\zeta_3 \times G(a_1, a_2, a_3, x)$, $\zeta_4 \times G(a_1, a_2, x)$, $\zeta_5 \times G(a_1, x)$, $\zeta_2 \zeta_3 \times G(a_1, x)$, $\pi^5 \times G(a_1, x)$, ζ_6 , ζ_3^2

$$z\bar{z} = p_2^2/p_1^2,$$

$$(1-z)(1-\bar{z}) = p_3^2/p_1^2$$

Full alphabet (from SOFIA)

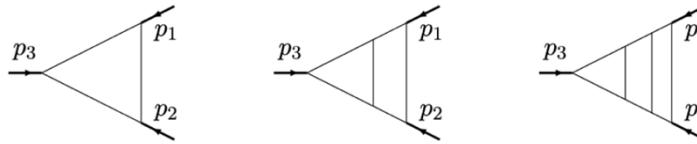
$$A_{1,2} = \left\{ z\bar{z}, (1-z)(1-\bar{z}), z-\bar{z}, \frac{\bar{z}}{z}, \frac{1-z}{1-\bar{z}} \right\}$$

Simplified alphabet (quicker for testing)

$$A_3^\star = \left\{ z\bar{z}, (1-z)(1-\bar{z}), \frac{\bar{z}}{z}, \frac{1-z}{1-\bar{z}} \right\}$$

Example 1: Triangles

Triangle ladder diagrams



Pick random points in (unphysical) Euclidean region $0 < z < \bar{z} < 1$

- Makes basis functions real
- Imaginary numbers are fine, just technically complicated python implementation

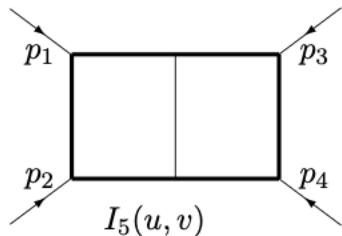
Diagram	AMFlow point time	Transcendental weights	# points sampled	Basis size	Reduction time	
One-loop	15.6 CPU-min	≤ 2	5	full(32)	<1s	3 digits
				simplified(26)	<1s	
				uniform(25)	<1s	
Two-loop	1.1 CPU-h	≤ 4	100	full(488)	9.6 min	20 digits
			100	simplified(393)	10.7 min	
			60	uniform(366)	3.5 min	
Three-loop	5.7 CPU-h	≤ 6	-	full(1373)	-	20 digits
			-	simplified(972)	-	
			200	uniform(806)	1.1 h	

Results with timing

- Rate limiting step is AMFlow (computation of full function)
- Could be sped up (additional points much faster than first)

Example 2: Double box

Loop



12 independent letters

$$\tilde{A} = \left\{ u, v, 1+u, 1+v, u+v, 1+u+v, \frac{\beta_u - 1}{\beta_u + 1}, \frac{\beta_v - 1}{\beta_v + 1}, \frac{\beta_{uv} - 1}{\beta_{uv} + 1}, \frac{\beta_{uv} - \beta_u}{\beta_{uv} + \beta_u}, \frac{\beta_{uv} - \beta_v}{\beta_{uv} + \beta_v}, \frac{\beta_{uv} - \beta_u \beta_v}{\beta_{uv} + \beta_u \beta_v} \right\}$$

with $\beta_u = \sqrt{1+u}$, $\beta_v = \sqrt{1+v}$ and $\beta_{uv} = \sqrt{1+u+v}$.

- $12^4 = 20,736$ weight-4 symbols + (?) lower weight terms
- square root letters are hard to integrate analytically

method 1:
rationalize the square roots

$$u = \frac{(1-w^2)(1-z^2)}{(w-z)^2} \quad \text{and} \quad v = \frac{4wz}{(w-z)^2},$$

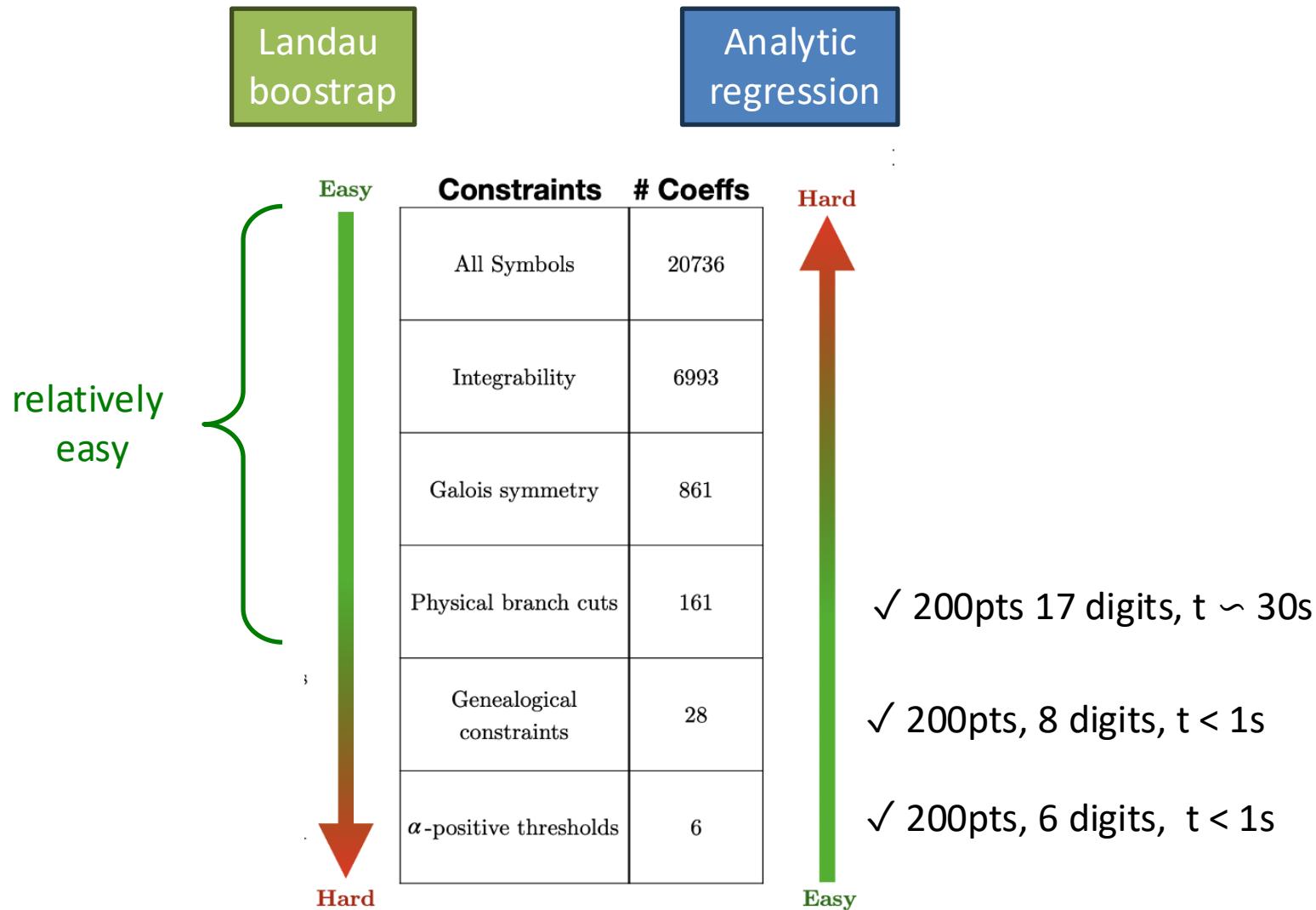
$$\tilde{A}_2 = \{w, z, 1 \pm w, 1 \pm z, w \pm z, 1 \pm wz, 1 \pm w \mp z + wz\}.$$

- Now symbols can be integrated analytically
 - Takes FiberSymbol hours to integrate
 - Result is hundreds or thousands of terms
- GiNaC can get numbers out, but very slow

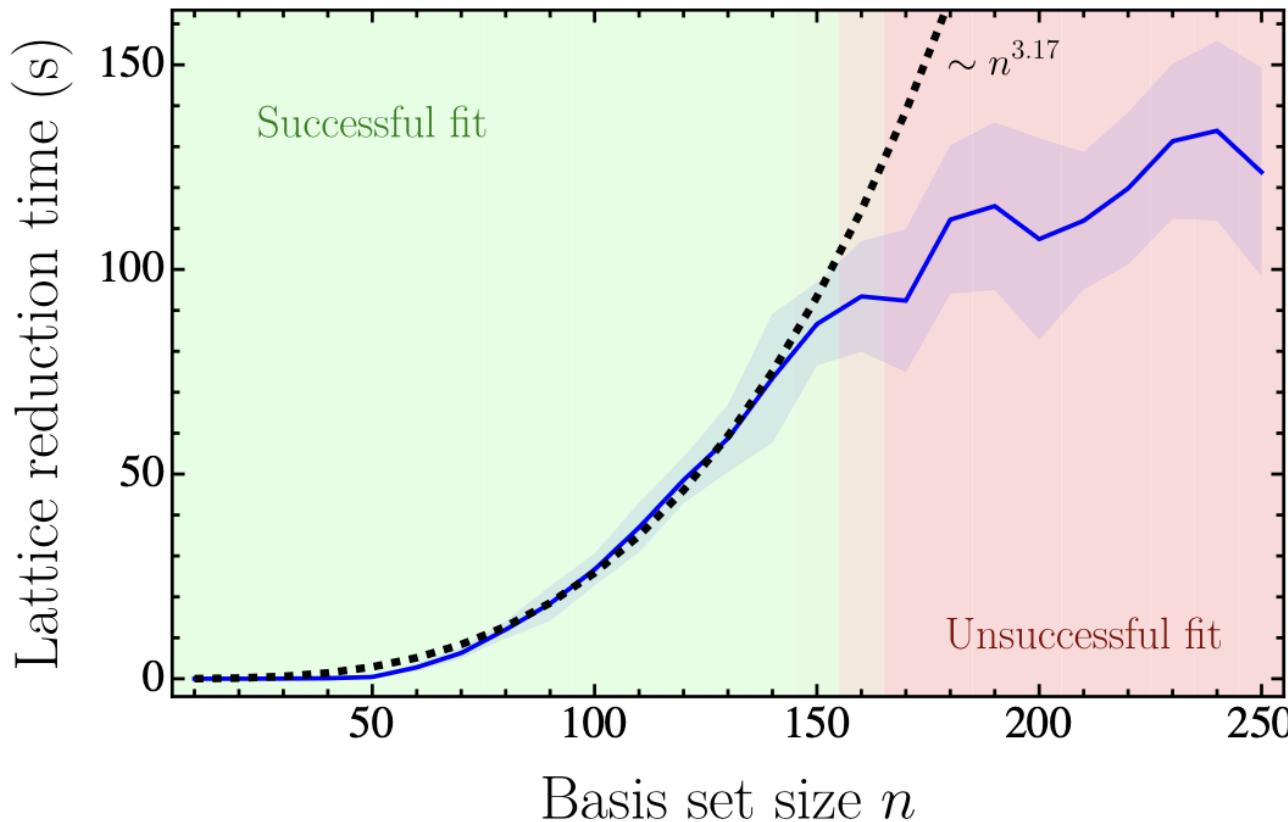
method 2:
numerically integrate along a contour

- integrate first and last symbol analytically
- need to be careful with branch cuts
 - euclidean region requires some thought
- using integrable contributions helps a lot

Example 2: Double box



Double box: limitations



Should work with more functions

- With our compute, hard to succeed with more than $n \lesssim 200$

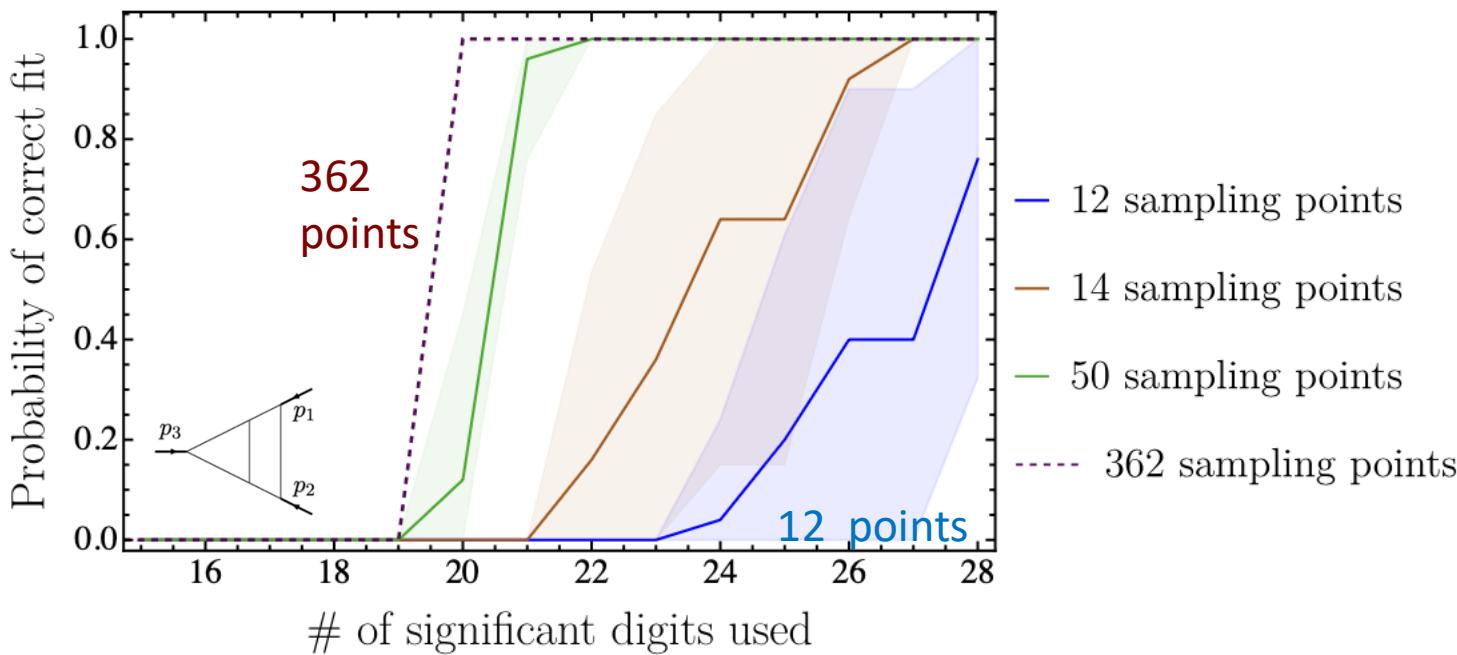
Choosing points

- Before we said you can get away with fewer digits if you use more points

$$d \lesssim \frac{nR}{p}$$

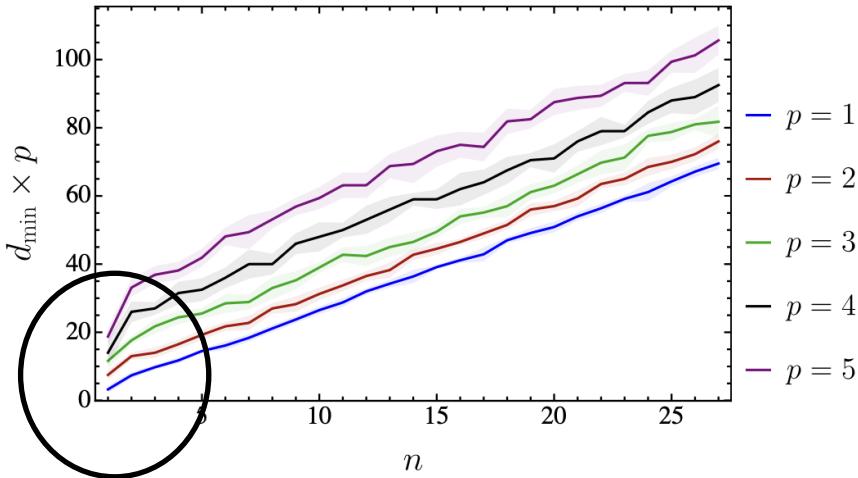
Annotations for the equation:

- # basis functions (green arrow)
- size of integers (orange arrow)
- number of points (purple arrow)
- digits of precision req'd (red arrow)



- Can never succeed below some digit lower bound
- Why did scaling fail?

Choosing points



offset near $d=0$

$$d_{\min} \approx R_{\text{eff}} \frac{n}{p} + d_0$$

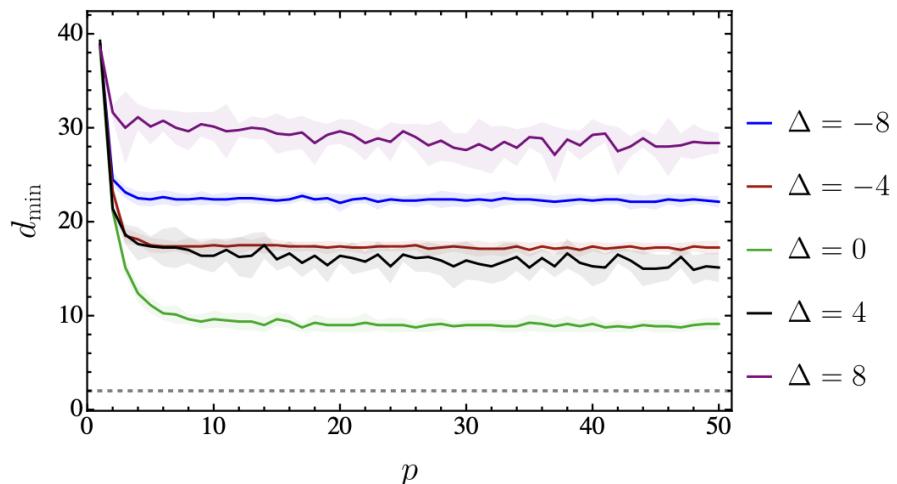
- will never work with 1 or 2 digits

$$f(p_j) = \{103.2, \quad \underbrace{2.5, \quad 2.3}_{\text{unrecoverable information}}\}$$

loss if we truncate to 2 digits

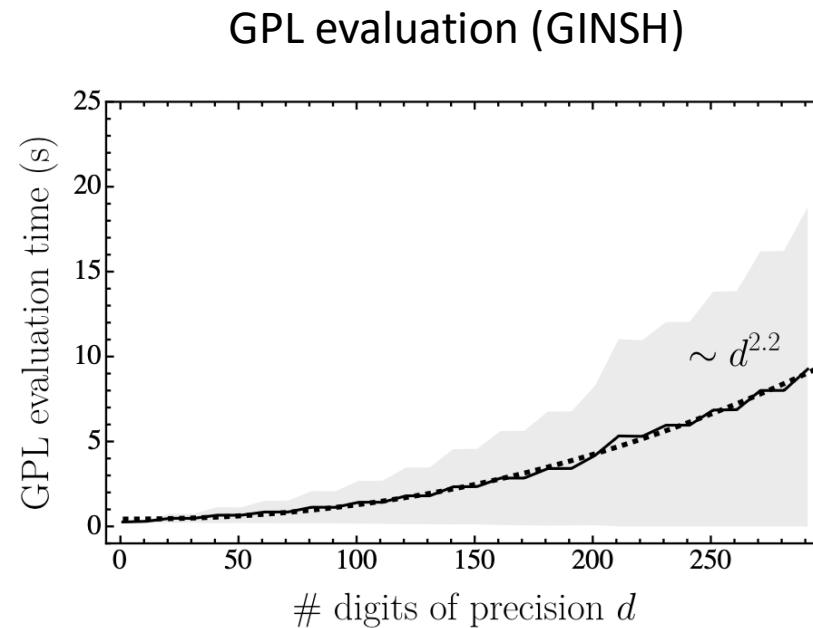
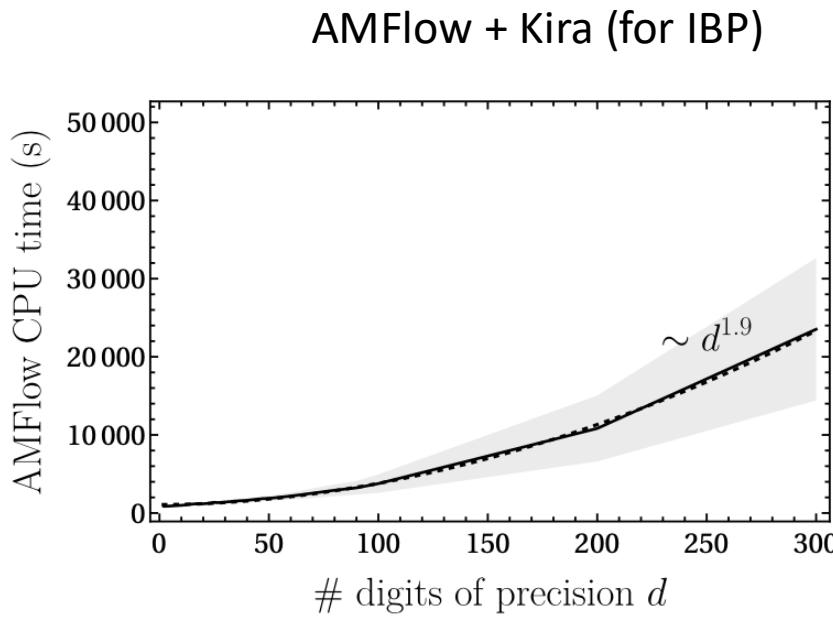
choose points in a range

$$\frac{1}{2} - 10^\Delta \leq x \leq \frac{1}{2} + 10^\Delta$$



- want to choose points
 - not too close (lose information)
 - not too far (need more digits)
 - Need 10-20 digits at least

Double box: timing



$$t(p, n)/\text{ns} \approx \underbrace{10^9 \cdot p \cdot d_{\min}^2(n)}_{\text{AMFLOW}} + \underbrace{10^4 \cdot n \cdot p \cdot d_{\min}^2(n)}_{\text{GINSH}} + \underbrace{p \cdot n^4}_{\text{fitting}},$$

- Trading digits for points makes scaling go from quadratic to linear!

Application to EECs

Toward the Analytic Bootstrap of Energy Correlators

Jianyu Gong^{1,2}, Andrzej Pokraka^{1,3}, Kai Yan^{1,2}, Xiaoyuan Zhang^{1,4,5}

2509.22782

Consider some types of energy-energy correlators (unequal energy weights)

$$\frac{d\sigma}{dx_{12} \cdots dx_{(N-1)N}} \equiv \sum_m \sum_{1 \leq i_1, \dots, i_N \leq m} \int d\sigma_m \times \prod_{1 \leq k \leq N} \frac{E_{i_k}}{Q} \prod_{1 \leq j < l \leq N} \delta \left(x_{jl} - \frac{1 - \cos \theta_{i_j i_l}}{2} \right).$$

Uses of lattice reduction

- Find linear dependence among basis functions

$$G(z) = D(z) + a_{0,11} g_1(z) + a_{0,12} g_2(z) + a_{1,12} g_3(z) + a_{2,12} g_4(z) + c_{1,12} g_5(z) + d_{2,12} g_6(z)$$

$$g_3(z) = \frac{1}{(z-1)^2 z^2 (\bar{z}-1)^2 \bar{z}^2} \left[z^4 \bar{z}^2 - z^4 \bar{z} + 8z^3 \bar{z}^3 - 14z^3 \bar{z}^2 + 8z^3 \bar{z} - z^3 + z^2 \bar{z}^4 - 14z^2 \bar{z}^3 + 24z^2 \bar{z}^2 - 14z^2 \bar{z} + z^2 + (2z^4 \bar{z}^3 - 3z^4 \bar{z}^2 + z^4 \bar{z} + 2z^3 \bar{z}^4 - 8z^3 \bar{z}^3 + 9z^3 \bar{z}^2 - 4z^3 \bar{z} + z^3 - 3z^2 \bar{z}^4 + 9z^2 \bar{z}^3 - 12z^2 \bar{z}^2 + 9z^2 \bar{z} - 3z^2 + z \bar{z}^4 - 4z \bar{z}^3 + 9z \bar{z}^2 - 8z \bar{z} + 2z + \bar{z}^3 - 3\bar{z}^2 + 2\bar{z}) \log((z-1)(\bar{z}-1)) + (-2z^4 \bar{z}^3 + 3z^4 \bar{z}^2 - z^4 \bar{z} - 2z^3 \bar{z}^4 + 8z^3 \bar{z}^3 - 9z^3 \bar{z}^2 + 2z^3 \bar{z} + 3z^2 \bar{z}^4 - 9z^2 \bar{z}^3 + 6z^2 \bar{z}^2 - z \bar{z}^4 + 2z \bar{z}^3) \log(z \bar{z}) - z \bar{z}^4 + 8z \bar{z}^3 - 14z \bar{z}^2 + 8z \bar{z} - \bar{z}^3 + \bar{z}^2 \right]$$

- Treat $\{g_i\}$ as transcendental basis and $a_{i,j}, b_{i,j}, c_{i,j}, d_{i,j}$ as coefficients to determine
- (1). Run lattice reduction among $\{g_i(z)\}$ to get a linear-independent basis $\{\tilde{g}_i(z)\}$
- (2). Run lattice reduction among both $G(z)$ and $\{\tilde{g}_i(z)\}$

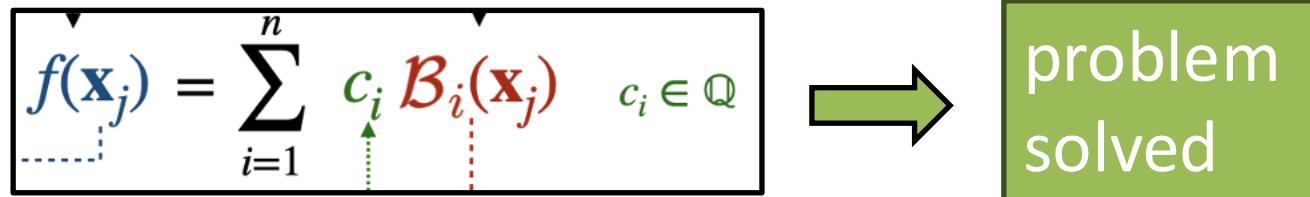
In this case, 2 numerical points and 13 digits fix all six parameters

Pros and cons: Landau bootstrap

- Can eliminate large swaths of symbols with physical constraints
- Don't need to do integrals
- Leads to new deep understanding in what amplitudes are
- Requires subtle understanding of singularities
 - Analytic structure of amplitudes
 - Branch points, euclidean regions
 - Algebraic geometry
- Rational prefactors not fixed by singularities alone
- Often still requires some integration at the end

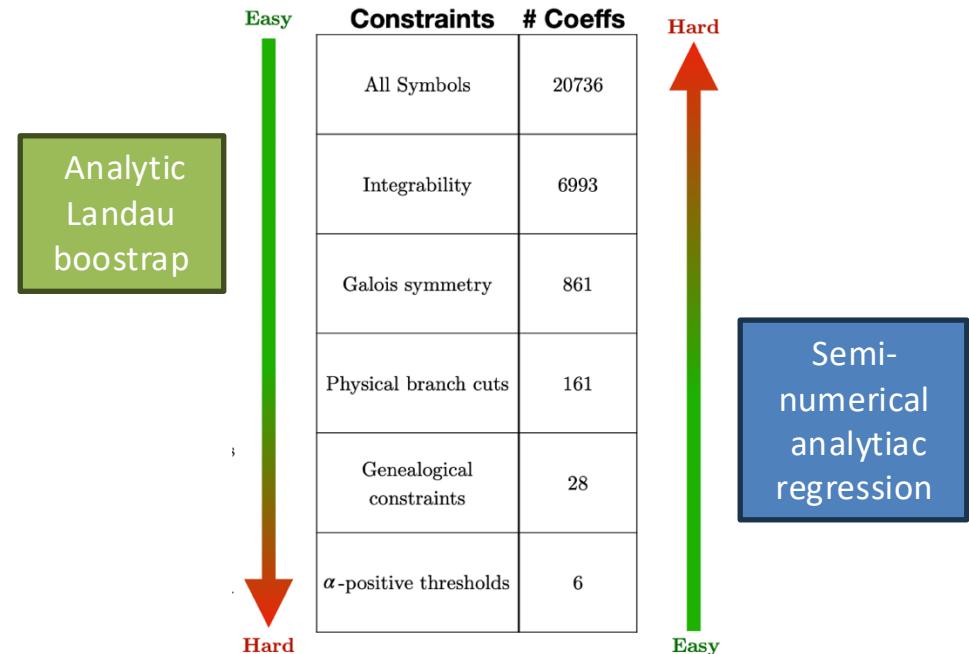
Pros and cons: Lattice reduction

- Easy to automate
- Can work for any functions
 - not just polylogs with symbols
 - elliptic polylogarithms? no problem!
 - cross sections, EECs, etc.
- Can find linear dependencies
- Can trade digits of accuracy for points
 - Scaling $t \sim n^2$ or worse to $t \sim n$
- Becomes computationally challenging above $n \sim 200$
 - AMFlow scales like $(\# \text{ digits})^2$
 - lattice reduction scales like $(\# \text{ constants})^4$
- Minimum number of digits needed (~ 5 or 6)



Conclusions

- Bootstrapping Feynman integrals is now possible in a semi-automated way
- Analytic and numerical methods are complementary



- Semi-numerical analytic regression could be widely valuable for many tasks